

# IA-32 Boot Sector Code

- **Boot Sector Code**
- **Boot Loader (Real & Protected Mode)**
- **Real-Mode Test Kernel**
- **Protected-Mode Test Kernel**

WEQAAR A. JANJUA  
<janjua@asu.edu>

Department of Computing Studies  
Arizona State University - Polytechnic

## **0. ABSTRACT**

The project is to write FAT-12 filesystem on a floppy disk and a bootloader that boots two different test kernels i.e. Real-mode and Protected-mode by first switching to the appropriate mode and handing the control over to the kernel. The kernels just print a string on bootup, they are not real kernels i.e. written just for the sole purpose of demonstrating boot loading process.

## **1. INTRODUCTION**

A computer system is a complex machinery, and the operating system is an elaborate tool that unrolls hardware complexities to end up showing a simple and standardized environment to the end user. When the power is turned on, however, the system software must work in a limited environment, and it must load the kernel using this scarce operating environment. This paper describes the booting process of IA-32 Platform.

## **2. THE COMPUTER AT POWER-ON**

In order to be able to do something with the computer when power is applied, things are arranged so that the processor begins execution from the system's firmware. The firmware is "unmovable software" found in ROM memory; some companies call it BIOS (Basic Input-Output System) to underline its software role, some call it PROM or "flash" to stress on its hardware implementation, while someone else calls it "console" to focus on user interaction.

The firmware usually checks that the hardware is correctly working, and retrieves part (or all) of the kernel from a storage medium and executes it. This first part of the kernel must load the rest of itself and initialize the whole system.

## **3. THE PC**

When the x86 processor is turned on in a personal computer, it is a 16-bit processor that only sees one Meg of RAM. This environment is known as "real mode", and is dictated by compatibility with older processors of the same family. It performs a POST(Power On Self Test) that initializes the chip set and checks that the computer is able to function correctly. Since the BIOS provide some basic hardware access it also initializes that and performs whatever internal house keeping that is necessary. One of the thing that is does is set up the BIOS Data Area.

When the BIOS is done starting up it loads the first sector of the floppy disk into memory at `0x0000:0x7c00`. That first sector is the *BOOT SECTOR*. The BIOS checks the format of the boot sector and will usually complain with some BIOS dependent message like "No system on disk" if it encounters an error.

### FORMAT OF THE BOOT SECTOR:

The boot sector is 512 bytes long and is the very first sector on a floppy disk. The first 3 bytes of the sector must be a jump or a short jump followed by a *NOP*.

```
jmp start_code  
; some code  
start_code:
```

OR

```
jmp short start_code  
nop ;required nop  
; some code  
start_code:
```

Some BIOS's reportedly check for the *NOP*. For FAT compatibility the next 59 bytes must contain the BIOS Parameter Block.

### THE BIOS PARAMETER BLOCK

Offset	Size	Description (default FAT12 1.44 Mb value)
0	8	Name of operating system
8	2	Bytes per sector (0x200)
10	1	Sectors per cluster (1)
11	2	Reserved sectors (1)
13	1	Number of FATS (2)
14	2	Root directory entries (0x00E0)
16	2	Total sectors (0x0B40)
18	1	Media Descriptor (0xF0)
19	2	Sectors per FAT (9)
21	2	Sector per track (0x12)
23	2	Number of heads (2)
25	4	Hidden sectors (0)
29	4	Total sectors huge (0)
33	1	Drive number (0)
34	1	Reserved
35	1	Signature (0x29)
36	4	Volume ID
40	11	Volume name <sup>1</sup>
51	8	File system type ('FAT12')

(All the strings must be padded with spaces.)

## FLOPPY DISK MEDIA DESCRIPTORS

Des crip tor	For ma t	Size	C yli n de rs	Hea ds	Se ct or s <sup>1</sup>	F A T si ze <sup>2</sup>	R o o t s i z e <sup>2</sup>
0xfe	160 Kb	5 1/4	40	1	8	?	?
0xfc	180 Kb	5 1/4	40	1	9	?	4
0xff	320 Kb	5 1/4	40	2	8	?	?
0xfd	360 Kb	5 1/4	40	2	9	4	7
0xf9	720 Kb	3 1/2	80	2	9	6	7
0xf9	1.2 Mb	5 1/4	80	2	15	14	1 4
0xf0	1.4 4 Mb	3 1/2	80	2	18	18	1 4
?	2.8 8 Mb	3 1/2	80	2	36	?	?

<sup>1</sup> Sectors per. cylinder

<sup>2</sup> Size in sectors

Then comes the code that will load an operating system from disk.

At the very end of the boot sector at *OFFSET 510* (just two bytes before the end) you must store the boot disk signature.

```
dw 0xAA55 ; boot disk signature
```

### FORMAT OF BOOT SECTOR

Offset (byte)	Size (bytes)	Description
0	3	Jump
3	59	BIOS Parameter Block
62	448	Your code
510	2	Signature

The BIOS leaves the number of the boot drive in the 'dl' register before transferring control to your boot sector by jumping to 0x0000:0x7C00. All other registers are undefined.

### MEMORY MAP:

This is a map of the first megabyte of memory right after the BIOS has transferred control to the boot sector code

Address	Size	Name
0x0000:0 x0000	1024 bytes	Interrupt Vector Table
0x0040:0 x0000	256 bytes	BIOS Data Area
0x0050:0 x0000	?	Free memory
0x07C0: 0x0000	512 bytes	Boot sector code
0x07E0:0 x0000	?	Free memory
0xA000:0 x0000	64 Kb	Graphics Video Memory
0xB000:0 x0000	32 Kb	Monochrome Text Video Memory
0xB800:0 x0000	32 Kb	Color Text Video Memory
0xC000: 0x0000	256 Kb <sup>1</sup>	ROM Code Memory
0xFFFF: 0x0000	16 bytes	More BIOS data

## POSSIBLE VALUES OF DL WHEN THE BOOT SECTOR CODE STARTS EXECUTING

Value of dl register	Corresponding drive
0x00	First floppy drive
0x01	Second floppy drive
0x80	First hard drive
0x81	Second hard drive

What is actually in the boot sector is the code to execute and possible some data too. Since the computer will attempt to execute the data in the boot sector it has to contain valid code.

The task of the boot sector is to prepare for and load the next step of the operating system. The simplest is to load an image from disk and transferring control to it immediately. But there is plenty of room left for doing more things in the boot sector. It could be entering protected mode (which is demonstrated practically in the project).

The first thing the boot sector should do after the jump is to initialize the data segment and set up a stack.

If you have some data in your boot sector, i.e. text to display on the screen, you have to initialize *ds* to a known value before using it to index the data with. Using the segment where the BIOS loaded the boot sector (*0x07C0*) is very convenient.

### INITIALIZING DS REGISTER:

```
mov ax, 0x07C0
mov ds, ax           ; setup ds register
```

### SETTING UP A STACK:

The stack can be put on any place, as long as it does not interfere with the location of the boot sector code or some other areas of reserved memory areas. You should also pay attention later when you load the kernel image or maybe relocate the code. I have chosen to place it at *0x9000:0x0000*. Without a stack it can be dangerous to call the BIOS, since you don't know whether it has its own stack or is using yours. If it is not set up probably it could possibly corrupt data or code.

```
mov ax, 0x9000
mov ss, ax           ; setup a stack
mov sp, 0x2000      ; 8 kb
```

## INITIALIZING DS REGISTER:

```
mov ax, 0x2000
mov ds, ax          ; setup ds to match new location
jmp 0x2000:0x0000  ; transfer control new location
```

The bootloader provides the user with two options:

- Load REAL-MODE KERNEL
- Load PROTECTED-MODE KERNEL

## TRANSFERRING CONTROL TO THE REAL MODE TEST KERNEL:

The bootloader loads the kernel at the very bottom of the memory. Therefore the boot sector code needs to be relocated. For the same reason I also relocate the BIOS Data Area by moving it to `0x7000:0x0000`. Then after the kernel has loaded it can extract the data that it needs from there. The next step is to reset the disk drive and read the kernel image into memory. The kernel is stored directly after the boot sector on the floppy disk, it fits into one sector.

The last step is to transfer the control to the test kernel. This is simply done using a `jmp`. But before that I set up the `ds` register so that the kernel doesn't have to do that. I prefer to enter the kernel in a known and stable state.

```
mov ax, 0x2000
mov ds, ax          ; setup ds to match new location
jmp 0x2000:0x0000  ; transfer control new location
```

## TRANSFERRING CONTROL TO THE PROTECTED MODE TEST KERNEL:

In order to switch from real-mode to protected mode:

- *CLI*: Disable interrupts, because the installed interrupts are all written for real mode and if an interrupt would occur after the mode switch, your system would probably reboot.
- Load the GDTR using *lgdt*, to set up the GDT.
- Execute a *mov CR0* instruction to set the PE bit of control register 0.
- Immediately after the *mov, cr0* instruction perform a far jump to clear the instruction prefetch queue, because it's still filled with real mode instructions and addresses.
- Reload all the segment registers except CS. (which is reloaded by the far jump)
- Load the Interrupt descriptor tables to make interrupts possible
- *STI*: Re-enable interrupts.
- Enable the A20 line to prevent memory wrap.
- Disable NMI (non-maskable interrupts)

### **ENABLE THE A20 ADDRESS LINE:**

In order to use the full amount of RAM plugged in your computer you have to enable the *a20* address line. As mentioned earlier enabling a line of the floppy controller can do this. Setting the appropriate bit can change the state of this line. This bit is the second bit of the AT keyboard controller output port (*port 064h*). So in theory we can enable the *a20* address line by simply setting this second bit.

### **DISABLING THE NMI:**

The NMI belongs to the interrupts issued by the hardware. But a NMI (Non Maskable Interrupt) is supplied to the processor directly and not via the 8259A PIC. The NMI usually reports a parity error when reading a byte from memory.

The problem is that you can't disable the NMI with the *CLI* instruction. However, there are times you have to disable it (e.g. when switching into protected mode). A register is provided for this purpose.

The NMI mask register allows to disable (or enable the NMI). This register is controlled by bit 7 of *port 0A0h* for the PC/XT and by bit 7 of *port 70h* for the AT and his successors. Note that in the AT the address register for the CMOS RAM and the real-time clock are also located at port address *70h*. You should take care of modifying bit 7 only.

## **4. WHAT IS PROTECTED MODE**

The 8088 CPU used in the original IBM PC was not very scalable. In particular, there was no easy way to access more than 1 megabyte of physical memory. To get around this while allowing backward compatibility, Intel designed the 80286 CPU with two modes of operation: real mode, in which the '286' acts like a fast 8088, and protected mode (now called 16-bit protected mode).

Protected mode allows programs to access more than 1 megabyte of physical memory, and protects against misuse of memory (i.e. programs can't execute a data segment, or write into a code segment). An improved version, 32-bit protected mode, first appeared on the '386 CPU'.



## DIFFERENCES BETWEEN REAL- AND PROTECTED MODES

	REAL MODE	16-BIT PROTECTED MODE	32-BIT PROTECTED MODE
Segment base address	20-bit (1M byte range) = $16^*$ segment register	24-bit (16M byte range), from descriptor	32-bit (4G byte range), from descriptor
Segment size (limit)	16-bit, 64K bytes (fixed)	16-bit, 1-64K bytes	20-bit, 1-1M bytes or 4K-4G bytes
Segment protection	no	yes	yes
Segment register	segment base adr / 16	selector	selector

### PROTECTED MODE AND SEGMENTED MEMORY:

The segments are still there, but in 32-bit protected mode, you can set the segment limit to 4G bytes. This is the maximum amount of physical memory addressable by a CPU with a 32-bit address bus. Limit-wise, the segment then "disappears" (though other protection mechanisms remain in effect). This reason alone makes 32-bit protected mode popular.

### DESCRIPTOR:

In real mode, there is little to know about the segments. Each is 64K bytes in size, and you can do with the segment what you wish: store data in it, put your stack there, or execute code stored in the segment. The base address of the segment is simply 16 times the value in one of the segment registers.

In protected mode, besides the segment base address, we also need the segment size (limit) and some flags indicating what the segment is used for. This information goes into an 8-byte data structure called a descriptor:

### CODE/DATA SEGMENT DESCRIPTOR

Lowest byte	Byte 1	Byte 2	Byte 3	Byte 4	Byte 5	Byte 6	Highest byte
Limit 7:0	Limit 15:8	Base 7:0	Base 15:8	Base 23:16	Access	Flags, Limit 19:16	Base 31:24

This is a 32-bit ('386) descriptor. 16-bit ('286) descriptors have to top two bytes (Limit 19:16, Flags, and Base 31:24) set to zero. The Access byte indicates segment usage (data segment, stack segment, code segment, etc.):

### ACCESS BYTE OF CODE/DATA SEGMENT DESCRIPTOR

Highest bit	Bits 6, 5	Bit 4	Bits 3	Bit 2	Bit 1	Lowest bit
Present	Privilege	1	Executable	Expansion direction/ conforming	Writable/ readable	Accessed

**Present bit:** Must be set to one to permit segment access.

**Privilege:** Zero is the highest level of privilege (Ring 0), three is the lowest (Ring 3).

**Executable bit:** If one, this is a code segment, otherwise it's a stack/data segment.

**Expansion direction (stack/data segment):** If one, segment grows downward, and offsets within the segment must be greater than the limit.

**Conforming (code segment):** Privilege-related.

**Writable (stack/data segment):** If one, segment can be written to.

**Readable (code segment):** If one, segment can be read from. (Code segments are not writable.)

**Accessed:** This bit is set whenever the segment is read from or written to.

The 4-bit Flags value is non-zero only for 32-bit segments:

#### FLAGS NYBBLE

Highest bit	Bit 6	Bit 5	Bit 4
Granularity	Default Size	0	0

The granularity bit indicates if the segment limit is in units of 4K byte pages (G=1) or if the limit is in units of bytes (G=0). For stack segments, the default Size bit is also known as the B (Big) bit, and controls whether 16- or 32-bit values are pushed and popped. For code segments, the D bit indicates whether instructions will operate on 16-bit (D=0) or 32-bit (D=1) quantities by default. To expand upon this: when the D bit is set, the segment is *USE32*, named after the assembler directive of the same name. The following sequence of hex bytes 'B8 90 90 90 90' will be treated by the CPU as a 32-bit instruction, and will disassemble as *mov eax, 90909090h*. In a 16-bit (*USE16*) code segment, the same sequence of bytes would be equivalent to

```
mov ax,9090h  
nop  
nop
```

Two special opcode bytes called the Operand Size Prefix and the Address Length Prefix reverse the sense of the D bit for the instruction destination and source, respectively. These prefixes affect only the instruction that immediately follows them.

Bit 4 of the Access byte is set to one for code or data/stack segments. If this bit is zero, you have a system segment. These come in several varieties:

**Task State Segment (TSS):** These are used to simplify multitasking. The '386 or higher CPU has four sub-types of TSS.

**Local Descriptor Table (LDT):** Tasks can store their own private descriptors here, instead of the GDT.

**Gates:** These control CPU transitions from one level of privilege to another. Gate descriptors have a different format than other descriptors:

### GATE DESCRIPTOR

<b>Lowest byte</b>	<b>Byte 1</b>	<b>Byte 2</b>	<b>Byte 3</b>	<b>Byte 4</b>	<b>Byte 5</b>	<b>Byte 6</b>	<b>Highest byte</b>
Offset 7:0	Offset 15:8	Selector 7:0	Selector 15:8	Word Count 4:0	Accesses	Offset 23:16	Offset 31:24

Note the Selector field. Gates work through indirection, and require a separate code or TSS descriptor to function.

### ACCESS BYTE OF SYSTEM SEGMENT DESCRIPTOR

<b>Highest bit</b>	<b>Bits 6, 5</b>	<b>Bit 4</b>	<b>Bits 3, 2, 1, 0</b>
Present	Privilege	0	Type

### SYSTEM SEGMENT TYPES

Type	Segment function	Type	Segment function
0	(Invalid)	8	(Invalid)
1	Available '286 TSS	9	Available '386 TSS
2	LDT	10	(Undefined, reserved)
3	Busy '286 TSS	11	Busy '386 TSS
4	'286 Call Gate	12	'386 Call Gate
5	Task Gate	13	(Undefined, reserved)
6	'286 Interrupt Gate	14	'386 Interrupt Gate
7	'286 Trap Gate	15	'386 Trap Gate

For now, TSSes, LDTs, and gates are the three main types of system segment.

### DESCRIPTORS

They are stored in a table in memory: the Global Descriptor Table (GDT), Interrupt Descriptor Table (IDT), or one of the Local Descriptor Tables.

The CPU contains three registers: GDTR, which must point to the GDT, IDTR, which must point to the IDT (if interrupts are used), and LDTR, which must point to the LDT (if the LDT is used). Each of these tables can hold up to 8192 descriptors.

**SELECTOR:**

In protected mode, the segment registers contain selectors, which index into one of the descriptor tables. Only the top 13 bits of the selector are used for this index. The next lower bit chooses between the GDT and LDT. The lowest two bits of the selector set a privilege value.

**HOW TO ENTER PROTECTED MODE**

Entering protected mode is actually rather simple.

**You must:**

- Create a valid Global Descriptor Table (GDT)
- (Optional) create a valid Interrupt Descriptor Table (IDT)
- Disable interrupts
- Point GDTR to your GDT
- (Optional) point IDTR to your IDT
- Set the PE bit in the MSW register
- Do a far jump (load both CS and IP/EIP) to enter protected mode (load CS with the code segment selector)
- Load the DS and SS registers with the data/stack segment selector
- Set up a pmode stack
- (Optional) enable interrupts

## SOURCE CODE SECTION

### BOOTSECTOR AND BOOTLOADER:

```
-----
; BOOT SECTOR CODE by Weqaar A. Janjua
-----

; to assemble: nasm bootsec.asm -f bin -o bootsec.bin
; to transfer to disk: partcopy bootsec.bin 0 200 -f0
[bits 16]
[org 0]
    jmp short start
    nop                ; required nop as some BIOS'es need it
-----
; BIOS PARAMETER BLOCK (definitions for protected mode)
-----
; FIELD SIZE (bytes)
    osname             db 'RAPTOR '           ; 8
    bytespersector     dw 0x200              ; 2
    sectorspercluster  db 1                 ; 1
    reservedsectors    dw 1                 ; 2
    numberoffats       db 2                 ; 1
    rootdirectoryentries dw 0x00E0 ;224      ; 2
    totalsectors       dw 0x0B40 ;2880      ; 2
    mediadescriptor    db 0xF0 ;1.44 MB     ; 1
    sectorsperfat      dw 2                 ; 2
    sectorspertrack    dw 0x12             ; 2
    numberofheads      dw 2                 ; 2
    hiddensectors      dd 0                 ; 4
    totalsectorshuge   dd 0                 ; 4
    drivenummer        db 0                 ; 1
    reserved           db 0                 ; 1
    signature          db 0x29             ; 4
    volumeid           dd 0                 ; 1
    volumename         db 'NONAME '        ; 1
    filesystemtype     db 'FAT12 '        ; 8
-----
; CODE
-----

; -----
; Functions used in the boot-loading process
; -----

start:
    cli                ; diable interrupts
    mov ax, 0x07C0
    mov ds, ax        ; setup ds register
    mov ax, 0x9000
    mov es, ax        ; setup a stack
    mov sp, 0x2000    ; 8 kb
    sti                ; enable interrupts
    mov [bootdrive], dl ; save boot drive
```

```

; relocate code
mov ax, 0x8000
mov es, ax
mov di, 0           ; destination address
mov si, 0           ; source address.
mov cx, 512        ; length is 512 bytes
cld                ; direction forward
rep movsb          ; move the boot sector
jmp 0x8000:relocation_ok ; transfer control to new location

```

relocation\_ok:

```

; relocate bios data area
mov ax, 0x7000
mov es, ax
mov di, 0           ; destination
mov ax, 0x0040
mov ds, ax
mov si, 0           ; source
mov cx, 256        ; length is 256 bytes
cld                ; set direction forward
rep movsb          ; move bios data area
mov ax, 0x8000
mov ds, ax         ; setup ds to match new location

```

call\_user:

```

mov si, ask_user
call bios_print_string
mov si, option1
call bios_print_string
mov si, option2
call bios_print_string
jmp o_a

```

o\_a:

```

mov ah, 0           ; wait for key
int 016h
cmp al, 'p'
je s_p
jne s_r
ret

```

s\_p: ; protected mode section

```

call bios_clear_screen
; load kernel image
call bios_reset_drive
jnc drive_ok
mov si, driveerr
call bios_print_string
call reboot

```

drive\_ok:

```

mov ax, 0x200
mov es, ax
mov bx, 0           ; kernel image destination
mov al, 1           ; read 1 sector
mov cl, 2           ; starting at sector 2

```

```
call bios_read_sectors
```

```
; ENABLE THE A20 LINE:
```

;In order to use the full amount of RAM plugged in your computer you have to enable the a20 addressline. This can be done by enabling a line of the floppy controller. The state of this line can be changed by setting the appropriate bit. This bit is the second bit of the AT keyboard controller output port (port 064h). So in theory we can enable the a20 address line by simply setting this second bit.

```
cli                ; disable interrupts
mov bl, 0xd0       ; read current status command
call kbd_send_ctrl_cmd
call kbd_read_data
or al, 2           ; set the a20 enable bit
push ax
mov bl, 0xd1       ; write current status command
call kbd_send_ctrl_cmd
pop bx
call kbd_write_data ; write the new status
mov bl, 0xd0       ; read current status command
call kbd_send_ctrl_cmd
call kbd_read_data ; read the current status
and al, 2
sti                ; enable interrupts
jnz a20_ok
mov si, a20err
call bios_print_string
call reboot
```

```
a20_ok:
```

```
; setup global descriptor table
```

```
mov ax, 0
mov es, ax
mov di, 0x800      ; destination
mov si, gdt        ; source
mov cx, 24        ; length
cld               ; forward direction
rep movsb         ; move gtd to its new location
lgdt [gdtptr]     ; load gdt register
```

```
; Disable ALL interrupts
```

```
cli                ; disable interrupts
mov al, 11111111b ; select to mask of all IRQs
out 0x21, al       ; write it to the PIC controller
```

```
;Disable NMI
```

```
in al, 0x70        ; read a value
or al, 10000000b   ; set the nmi disable bit
out 0x70, al       ; write it back again
```

```
; Enter protected mode
```

```
mov eax, cr0
or al, 1           ; set protected mode bit
mov cr0, eax
```



```
; Transfer control to kernel
mov ax, 0x10
mov ds, ax          ; load global data selector into ds
jmp 0x08:0x2000    ; transfer control to test kernel
```

```
ret
```

```
-----
; start of realmode kernel proc
-----
```

```
s_r:                ; real-mode section
```

```
    call bios_clear_screen
    call bios_reset_drive
    jnc reset_ok
    mov si,driveerr
    call bios_print_string
    call reboot
```

```
reset_ok:
```

```
    mov al, 1      ; sector count
    mov cl, 3      ; start sector
    mov ax, 0x2000
    mov es, ax
    mov bx, 0      ; kernel image destination
    call bios_read_sectors
    jmp 0x2000:0x0000 ; transfer control to kernel
    call call_user
```

```
ret
```

```
-----
; END of realmode kernel proc
-----
```

```
-----
; Functions
-----
```

```
bios_print_string:
```

```
; input : ds:si points to zero terminated string
    cld                ; direction forward
    lodsb              ; get next character
    cmp al, 0
    jz bios_print_string_done
    mov ah, 0x0E        ; write character as tty function
    int 0x10            ; call bios video services
    jmp bios_print_string
```

```
bios_print_string_done:
```

```
    ret
```

```
bios_clear_screen:
```

```
    mov al, 3          ; select video mode 3 - color text
    mov ah, 0          ; set video mode function
    int 0x10            ; call bios video services
    ret
```

```

reboot:
    mov si, pm
    call bios_print_string
    mov ah, 0           ; read keypress function
    int 0x16           ; call bios keyboard services
    jmp 0xFFFF:0x0000

bios_reset_drive:
    mov ah, 0           ; reset drive function
    int 0x13           ; call bios disk i/o
    ret

bios_read_sectors
; input : es:bx = address of destination
;     al = sector count
;     cl = sector start number
    mov ah, 0x02       ; read sectors function
    mov ch, 0           ; cylinder 0
    mov dl, [bootdrive] ; drive number
    mov dh, 0           ; head number
    int 0x13           ; call bios disk i/o
    jc bios_read_sectors
    ret

kbd_wait_cmd:
    in al, 0x64         ; read the controller status port
    and al, 2           ; check if the controller is ready
    jnz kbd_wait_cmd   ; to accept the next command (or
    ret                 ; piece of data)

kbd_wait_data:
    in al, 0x64         ; read the controller status port
    and al, 1           ; check if the data is ready
    jz kbd_wait_data
    ret

kbd_send_ctrl_cmd:
; input : bl = command
    call kbd_wait_cmd
    mov al, bl
    out 0x64, al        ; send the command to the control
    ret                 ; register

kbd_read_data:
; output : al = data
    call kbd_wait_data
    in al, 0x60         ; read data from input/output port
    ret

kbd_write_data:
; input bl = data
    call kbd_wait_cmd
    mov al, bl
    out 0x60, al        ; write data to input/output port
    ret

```

```

;-----
; data
;-----

; messages (with carriage return and line feed and zero terminated)
    bm_p    db  'PROTECTED', 0
    bm_r    db  'REAL', 0
    pm      db  'rb', 0 ; reboot message
    driveerr db  'E', 0 ; DRIVE ERROR
    a20err  db  'E', 0 ; A20 ERROR
    ask_user db  'SELECT:',13,10,0
    option1 db  'p',13,10,0
    option2 db  'r',13,10,0
    bootdrive db  0

; global descriptor table
; null selector (required)
    gdt dw 0, 0, 0, 0

; kernel code selector
    dw 0xffff ; segment limit (4 gb total)
    dw 0 ; base address (bits 0-15)
    db 0 ; base address (bits 16-24)
    db 10011000b ; dpl 0, code (execute only)
    db 11001111b ; granlurarity (4k), 32-bit, limit high nibble = f
    db 0 ; base address (bits 24-32)

; kernel data selector
    dw 0xffff ; segment limit (4 gb total)
    dw 0 ; base address (bits 0-15)
    db 0 ; base address (bits 16-24)
    db 10010010b ; dpl 0, data (read/write)
    db 11001111b ; granlurarity (4k), 32-bit, limit high nibble = f
    db 0 ; base address (bits 24-32)

    gdtpr dw 0x7ff ; limit (256 slots)
    dd 0x800 ; base (physical address)

;-----
; signature
;-----

    times 510-($-$$) db 0 ; padding - fill the empty space with 512 bytes !!
    dw 0xAA55 ; boot signature

;←-----END OF BOOTLOADER & BOOTSECTOR CODE----→

```

## REAL-MODE TEST KERNEL:

[bits 16]

start:

```

    call ClrScr
    mov ah,13h
    mov al,3 ; write mode (advance cursor, ASCII+attribute string)

```

```
mov bh,0 ; video page
mov cx,7 ; string length
mov dh,1 ; starting row
mov dl,1 ; starting col
push cs
pop es
mov bp,kernelmsg
int 10h
call reboot
```

```
-----
; functions
;-----
```

gotoxy:

```
mov ah,02h ; select video service 2 (position cursor)
mov bh,0 ; stay with video page 0
int 10h
```

ret

ClrScr:

```
pusha
mov cx,0
mov dx,LRXY
```

ClrWin:

```
mov al,0
```

ScrIWin:

```
mov bh,07h
```

video6:

```
mov ah,06h
int 10h
popa
```

ret

bios\_print\_string:

; input : ds:si points to zero terminated string

```
cld ; direction forward
lods b ; get next character
cmp al,0
jz bios_print_string_done
mov bh,0 ; setting video page (0)
mov bl,14h ; blue background and red font
mov ah,0x0E ; write character as tty function
int 0x10 ; call bios video services
jmp bios_print_string
```

bios\_print\_string\_done:

```
ret
```

reboot:

```
mov ah,13h
mov al,3 ; write mode (advance cursor, ASCII+attribute string)
mov bh,0 ; video page
```

```

        mov bl,02          ; attribute (black on green)
        mov cx,23         ; string length
        mov dh,2         ; starting row
        mov dl,2         ; starting col
        push cs
        pop es
        mov bp,presskeymsg
        int 10h
        mov ah, 0         ; read keypress function
        int 0x16         ; call bios keyboard services
        call do_reset

ret

do_reset:
        jmp 0xFFFF:0x0000

ret
;-----
; data
;-----

kernelmsg    db 'S', 01, 'U', 02, 'C', 03, 'C', 04, 'E', 05, 'S', 06, 'S', 07, 13, 10
presskeymsg  db 'P',01,'R',02,'E',03,'S',04,'S',05,' ',06,'A',07,'N',08,'Y',09,'
              ',10,'K',11,'E',12,'Y',13,',',14,'T',15,'O',16,'
              ',17,'R',18,'E',19,'B',20,'O',21,'O',22,'T',13,10
LRXY        dw          184Fh
times 512-($-$$) db 0    ; padding

;←-----END OF REAL-MODE TEST KERNEL CODE-----→

```

## PROTECTED-MODE TEST KERNEL:

```

[bits 32]
[org 0x2000]
        mov ax, 0x10
        mov ds, ax
        mov es, ax
        mov esi, kernelmsg
        call pmode_print_string
        mov esi, presskeymsg
        call pmode_print_string

dummy:
        jmp dummy
;-----
; function
;-----

        xposition db 0
        yposition db 1

pmode_print_character:
; input al : character
; ah : attribute
        pushad          ; save registers
        cmp al, 10     ; line feed

```

```
    jnz not_line_feed
    add byte [yposition], 1
    jmp pmode_print_character_done
```

```
not_line_feed:
    cmp al, 13                ; carriage return
    jnz not_carriage_return
    mov byte [xposition], 0
    jmp pmode_print_character_done
```

```
not_carriage_return:
    mov ecx, eax              ; save character and attribute
    mov ebx, 0
    mov bl, [xposition]
    shl bl, 1                ; calculate x offset
    mov eax, 0
    mov al, [yposition]
    mov edx, 160
    mul edx                   ; calculate y offset
    mov edi, 0xb8000         ; start of video memory
    add edi, eax              ; add y offset
    add edi, ebx              ; add x offset
    mov ax, cx                ; restore character and attribute
    cld                       ; forward direction
    stosw                     ; write character and attribute
    add byte [xposition], 1
```

```
pmode_print_character_done:
    call hardware_move_cursor
    popad                     ; restore registers
    ret
```

```
pmode_print_string:
; input ds:esi = points to zero terminated string
    lodsb
    cmp al, 0
    jz pmode_print_string_done
    mov ah, 0x0F              ; white text, black background
    call pmode_print_character
    jmp pmode_print_string
```

```
pmode_print_string_done:
    ret
```

```
hardware_move_cursor:
    pushad                    ; save registers
    mov ebx, 0
    mov bl, [xposition]       ; get x offset
    mov eax, 0
    mov al, [yposition]
    mov edx, 80
    mul edx                   ; calculate y offset
    add ebx, eax               ; calculate index
    ; select to write low byte of index
    mov al, 0xf
    mov dx, 0x03d4
```

```
out dx, al
; write it
mov al, bl
mov dx, 0x03d5
out dx, al
; select to write high byte of index
mov al, 0xe
mov dx, 0x03d4
out dx, al
; write it
mov al, bh
mov dx, 0x03d5
out dx, al
popad                ; restore registers
ret
```

```
;-----
; data
;-----
```

```
kernelmsg    db 'Protected mode test kernel loaded successfully', 13, 10, 0
presskeymsg  db 'Please remove disk and reboot', 0
times 512-($-$$) db 0      ; padding
```

```
;←-----END OF PMODE TEST KERNEL CODE-----→
```

```
;-----
; END of code
;-----
```